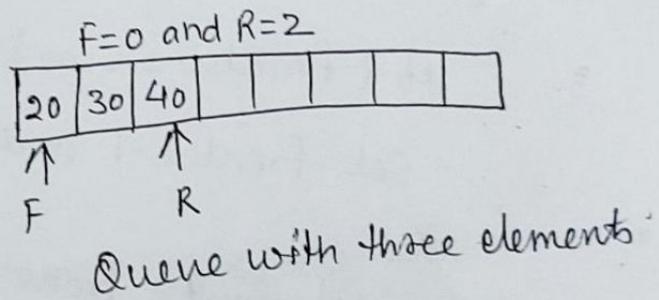
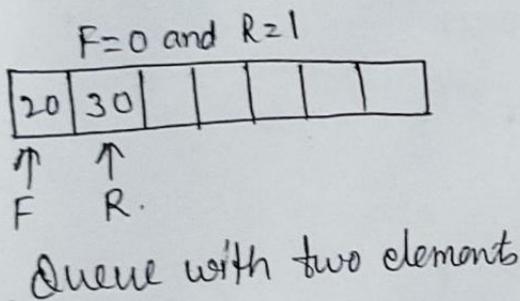
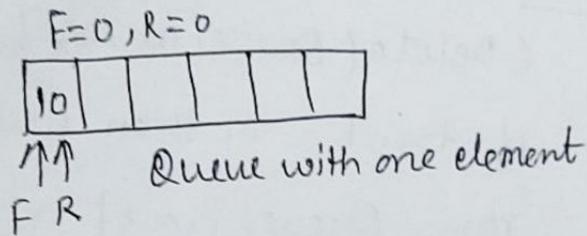
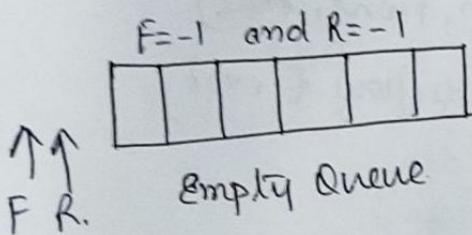


Queue → Queue is logically a first in first out types of list.

→ Queue is a non primitive linear data structure.

→ Queue is a homogeneous collection of elements in which new elements are added at one end called rear and existing elements are deleted from other end called the front.



Note - if $F = R$ then only one element is a Queue.

No. of elements in a Queue = $R - F + 1$

Implementation of Queue →

- 1) Using array
- 2) using linked list

1) Using array → $QInsert(Queue[maxsize], ITEM, front, rear)$
 Queue[maxsize] is the array, item is inserted or deleted in a Queue. front and rear is used for ~~insertion~~ deletion and insertion respectively.

②

1. if $\text{rear} == \text{maxsize} - 1$ then write overflow and exit.
2. if $\text{rear} == -1$ then set $\text{rear} = 0$ and $\text{front} = 0$
else set $\text{rear} = \text{rear} + 1$
3. $\text{Queue}[\text{rear}] = \text{item}$
4. Exit

2) Deletion -

QDelete (Queue[maxsize], item, front, rear)

1. if $\text{front} == -1$ then write underflow & exit
2. $\text{item} = \text{Queue}[\text{front}]$
3. if ($\text{front} == \text{rear}$) then
Set $\text{front} = -1$ and $\text{rear} = -1$
else
set $\text{front} = \text{front} + 1$
4. Exit

3) Traverse -

Qtraverse (Queue[maxsize], front, rear)

1. if $\text{rear} == -1$ then write Queue is empty & exit.
2. ~~for~~ $i = \text{front}$
3. repeat step 4 & 5 until $i < \text{rear}$
4. Display $\text{Queue}[i]$
5. $i = i + 1$
[END of Loop]
6. Exit

C function → Insertion

③

```
#define maxsize 30  
int queue[maxsize], front=-1, rear=-1;
```

```
void insert( )
```

```
{
```

```
    int item;
```

```
    if (front == (rear+1) % maxsize)  
    if (rear == maxsize-1) {  
    if (rear == maxsize-1)
```

```
        printf("Overflow");
```

```
        exit(1);
```

```
    else
```

```
    {
```

```
        if (rear == -1)
```

```
        {
```

```
            front = 0;
```

```
            rear = 0;
```

```
        }
```

```
    else
```

```
        rear = (rear+1) % maxsize;
```

```
        printf("Enter item");
```

```
        scanf("%d", &item);
```

```
        queue[rear] = item;
```

```
    }
```

```
}
```

void deletion ()

{

if (front = -1)

printf (" Underflow ");

else

{

if (front == rear)

{

front = -1;

rear = -1;

}

else

front = (front + 1); ~~rear = rear;~~

}

}

void display ()

{

int i;

if (rear == -1)

printf (" queue is empty ");

else

{

for (i = front; i <= rear; ~~++i~~ i = (i + 1) ~~++i~~)

printf (" %d ", queue [i]);

~~printf (" %d ", queue [i]);~~

}

}

void main ()

{

insert ();

deletion ();

insert ();

display ();

}

2) Linked List Implementation

Algorithm →

Insert (Queue, front, rear, info, next, item, new)

1. new = AVAIL
2. if (new == NULL) write overflow & exit.
3. Info[new] = item
4. LINK[new] = NULL
5. if (front == NULL) then
front = new and rear = new
6. else
~~rear → next =~~
LINK[rear] = new
Rear = new
7. Exit

Deletion (Queue, front, rear, PTR, next, INFO)

1. if (front == NULL)
write ~~over~~ underflow and exit
2. PTR = front
3. Display INFO[PTR] is deleted
4. if (~~front~~ → LINK[front] = NULL) then
front = NULL and rear = NULL & exit.
5. front = LINK[front]
6. free (PTR)
7. exit.

⑥ Traverse()

1. if (front == NULL) then write Queue is empty first.
2. ptr = front
3. repeat step 4 & 5 till ptr != NULL
4. display ptr->info
5. ptr = ptr->next
6. exit.

C function → ~~void insert~~

~~void insert~~

```
struct node
{
    int info;
    struct node *next;
} *front = NULL, *rear = NULL;
```

void insert()

{

struct node *new;

new = (struct node *) malloc(sizeof(struct node));

if (new == NULL)

printf("overflow");

else

{

printf("Enter item");

scanf("%d", &new->info);

new->next = NULL;

(17) (7)

```

if ( front == NULL )
{
    front = new;
    rear = new;
}
else
{
    rear->next = new;
    rear = rear;
}
}

```

```

void deletion()
{
    struct node *p;
    if ( front == NULL )
        printf( "Underflow" );
    else
    {
        p = front;
        printf( "%d is deleted", p->info );
        if ( p->next == NULL ) or if ( front == rear )
            front = rear = NULL;
        else
            front = front->next;
        free ( p );
    }
}
}

```

```

⑧ void traverse()
{

```

```

    struct node *p;

```

```

    if (Rear == NULL)

```

```

        printf("Queue is empty");

```

```

    else

```

```

        for (p = front; p != NULL; p = p->next)

```

```

            printf("%d", p->info);

```

```

    }

```

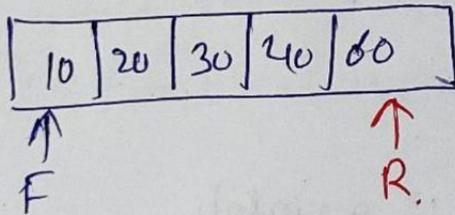
Limitation of ~~Circular~~ ^{Simple} Queue

❖ Overflow condition

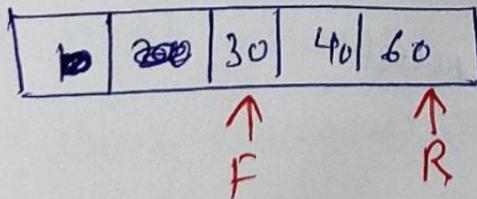
```

if (rear == maxsize - 1)

```



Deleted ~~two~~ Two items



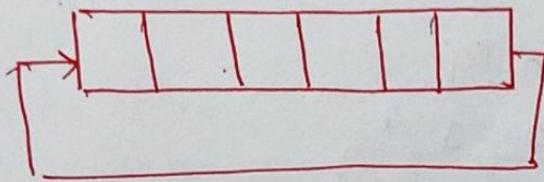
In above $R = \text{maxsize} - 1$, so no item can be inserted even there is a space in queue for two items.

Circular Queue →

9

problem of simple queue can be solved by circular queue.

In a circular queue the insertion of new element is done at first location of Queue if the last location of the Queue is full.



Overflow condition (Array implementation)

if (front = (rear + 1) % maxsize)

printf("Overflow");

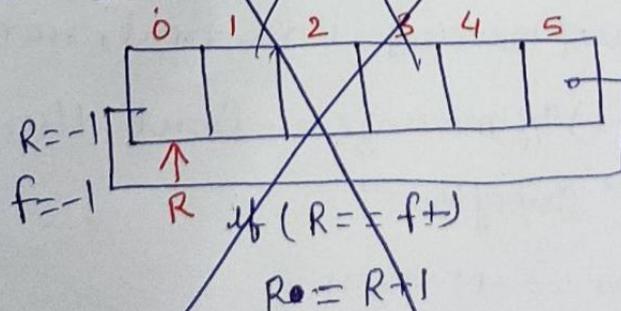
Underflow condition

front = -1 or rear = -1

operations on circular Queue → (Array implementation)

1) Insert →

concept → case 1 → Queue is empty

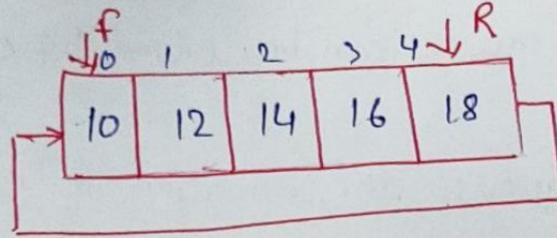


⑩

Insert →

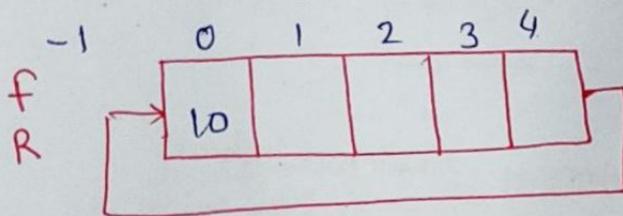
Case 1:- Queue is full

$M=5$ (max-size)



if $((R+1) \% \text{maxsize} == \text{front})$
 printf("Overflow")

Case 2:- Queue is empty



~~if (R == -1)~~

item = 10

if $(R == -1)$

f = R = 0

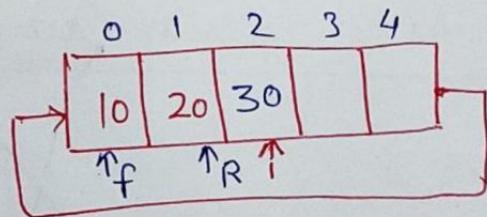
cq[R] = item

Case 3:-

Queue is neither empty nor full

$R = (R+1) \% \text{maxsize}$ item = 30

cq[R] = item



Algorithm →

Qinsert (Queue[maxsize], item, front, rear)

1. if $(\text{rear} + 1) \% \text{maxsize} == \text{front}$ then
write "Overflow"

2. if $(\text{Rear} == -1)$ then

Set ~~front~~ front = 0, rear = 0
else

Set rear = $(\text{rear} + 1) \% \text{maxsize}$

3. Queue[rear] = item

4. Exit.

(11)

C function →

```
#define maxsize 30
int eq[maxsize], front=-1, rear=-1 // global
                                   declaration

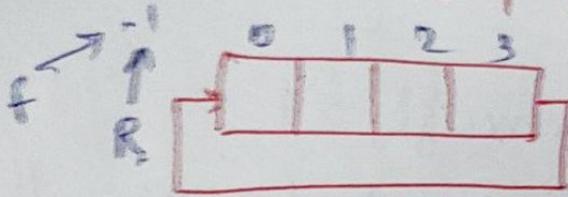
void insert()
{
    int item;

    if ((rear+1)%maxsize == front)
    {
        printf("Overflow");
        exit(0);
    }
    else
    {
        if (rear == -1)
            front = rear = 0;
        else
            rear = (rear+1)%maxsize;
        printf("Enter item");
        scanf("%d", &item);
        eq[rear] = item;
    }
}
```

2) Deletion →

(12)

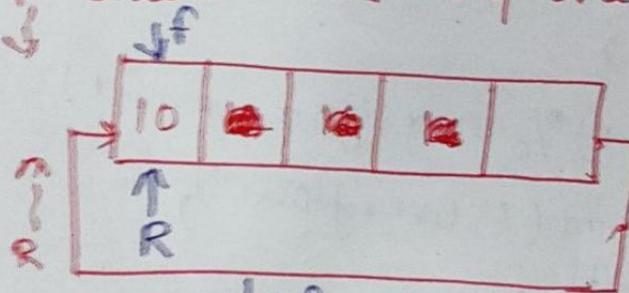
Case 1:- Queue is empty



if (front == -1)

printf("Underflow");

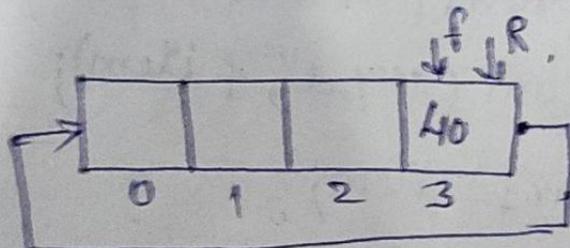
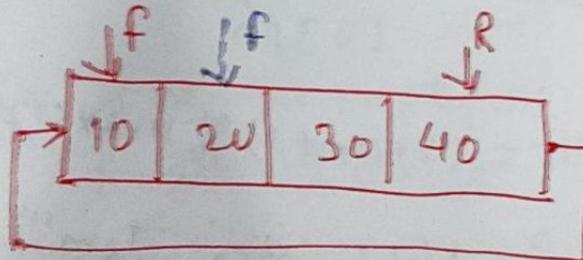
Case 2: Queue has only one item



if (front == rear)

front = rear = -1

Case 3: Other wise



front = (front + 1) % maxsize

Algorithm →

(13)

~~@@~~ Circular-Queue-delete ()

1. if (front == -1) then
write "Underflow" and exit.
2. item = Queue[front]
3. if (front == rear) then
front = rear = -1
else
front = (front + 1) % maxsize
[End of if]
4. write item "is deleted"
5. Exit.

C function →

```
void delete() ( )  
{  
    int item;  
    if if (front == -1) then  
    {  
        printf("Underflow");  
        exit(0);  
    }  
    item = cq[front];  
    if (front == rear)  
        front = rear = -1;  
    else  
        front = (front + 1) % maxsize;  
    printf("%d is deleted", item);  
}
```

Traverse → Concept →

(14) C++

Case 1:- Queue is empty

```
if (front == -1)
    printf("Queue is empty");
```

Case 2:- Queue is not empty

```
for (i = front; i != rear; i = (i+1) % maxsize)
    printf("%d", cq[i]);
printf("%d", cq[i]);
```

Algorithm →

Traverse()

1. if (front == -1)
write "Queue is empty & exit."
2. ~~for~~ i = front
3. repeat step 4 & 5 ~~while (i != rear)~~ while (i != rear)
4. write Queue[i]
5. i = (i+1) mod maxsize
6. write Queue[i]
7. Exit

C: function →

```
void traverse()
{
    if (front == -1)
    {
        printf("Queue is empty");
        exit(0);
    }
    for (i = front; i != rear; i = (i+1) % maxsize)
        printf("%d", cq[i]);
    printf("%d", cq[i]);
}
```

Circular Queue (Linked List implementation) →

```
struct node
{
    int info;
    struct node *next;
} *front = NULL, *rear = NULL;
```

1) insertion → concept

case 1 → Over flow condition
new = malloc()

```
if (new == NULL)
    printf("overflow");
```

Case 2 →

Queue is empty

(16)

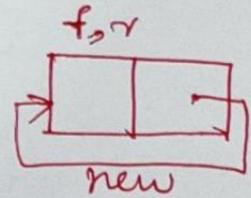
if (~~front~~ rear == NULL)

f, r = NULL
↓

front = new;

rear = new;

rear → next = front;

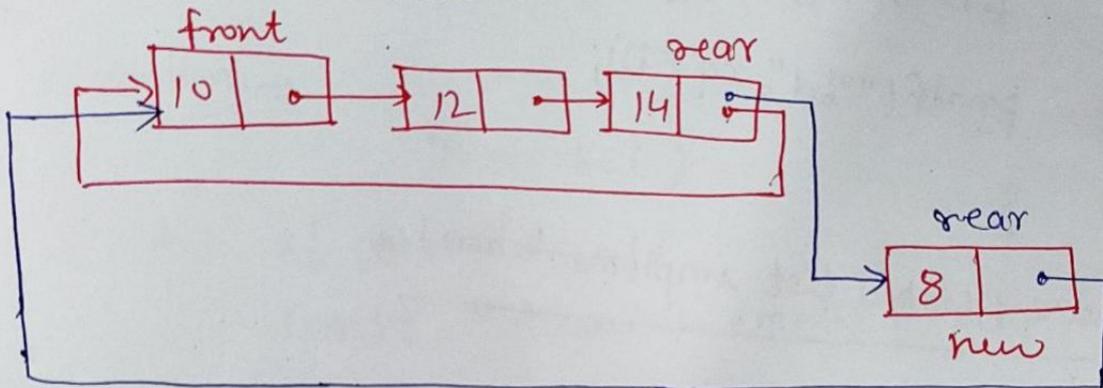


Case 3 -

rear → next = new

rear = new

rear → next = front



Algorithm →

insert ()

1. new = AVAIL
2. if (new == NULL)
write "Overflow" & exit.
3. if (rear == NULL) then
front = rear = new
~~rear~~ → LINK[rear] = front
else
LINK[rear] = new
rear = new
LINK[rear] = front.
4. Exit.

C function →

(17)

```
void insert()
{
    struct node *new;
    new = (struct node *) malloc (sizeof (struct node));
    if (new == NULL)
    {
        printf ("overflow");
        exit(0);
    }
    if (rear == NULL)
    {
        front = rear = new;
        rear->next = front;
    }
    else
    {
        rear->next = new;
        rear = new;
        rear->next = front;
    }
}
```

2) Deletion → concept:-

Case 1 →

Queue is empty

if (front == NULL)

printf ("underflow");

Case 2 →

Queue contains only one node

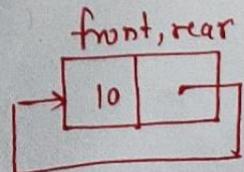
if (front == rear)

{

p = front

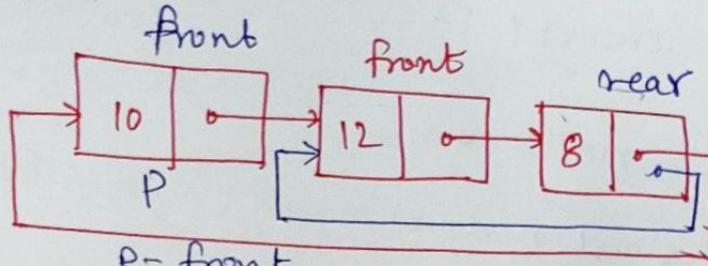
front = rear = NULL
free(p) }

F R
↓ ↓
NULL



case 3 → Queue contains more than one node.

(18)



$P = \text{front}$
 $\text{front} = \text{front} \rightarrow \text{next}$
 $\text{rear} \rightarrow \text{next} = \text{front}$
 $\text{free}(P)$

Algorithm →

deletion()

1. if (front == NULL) then
write "Queue is underflow" & exit.
2. $P = \text{front}$
3. if (front == rear) then
 $\text{front} = \text{rear} = \text{NULL}$
 $\text{free}(P)$
else
 $\text{front} = \text{front} \rightarrow \text{next};$
 $\text{rear} \rightarrow \text{next} = \text{front};$
 $\text{free}(P);$
4. exit

if (front == rear) then
 $\text{front} = \text{rear} = \text{NULL}$
 $\text{free}(P)$
else
 $\text{front} = \text{LINK}[\text{front}]$
 $\text{LINK}[\text{rear}] = \text{front}$
 $\text{free}(P)$
4. Exit.

C function →

```
void deletion()
{
    struct node *P;
    if (front == NULL)
    {
        printf("Queue is underflow");
        exit(0);
    }
}
```

(19)

```

P = front;
if (front == rear) then
{
    front = rear = NULL;
    free(P);
}
else
{
    front = front->next;
    rear->next = front;
    free(P);
}
}

```

3) Traverse or display-

Concept →

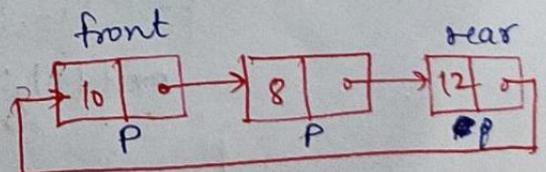
Case 1: Queue is empty
 if (front == NULL)
 printf("Queue is empty");

Case 2:- Queue is not empty

```

P = front;
do
{
    printf("%d", P->info);
    P = P->next;
} while (P != front) while (P != front)

```



Algorithm

(20)

Traverse()

1. if (front == NULL) then
write "Queue is empty" & exit.
2. P = front
3. ~~do~~ write INFO[P]
4. P = LINK[P]
5. while (P != front) repeat steps 5 & 6
6. write INFO[P]
6. P = LINK[P]
- [END of while]
7. Exit.

C function →

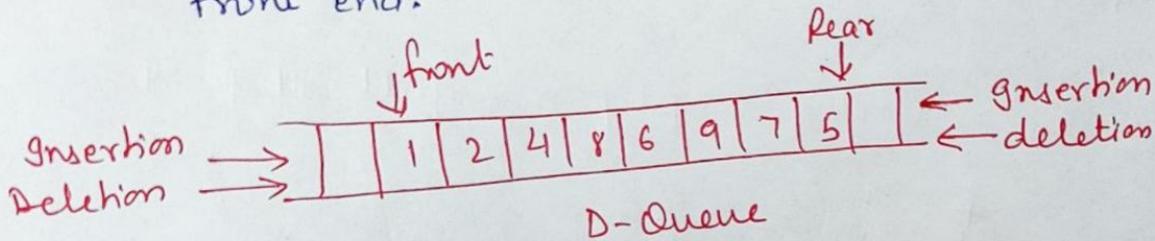
```
void traverse()
{
    struct node *p;
    if (front == NULL)
    {
        printf("Queue is empty");
        exit(0);
    }
    do P = front;
    do
    {
        printf("%d", P->info);
        P = P->next;
    } while (P != front);
}
```

}

Double ended Queue \rightarrow (D-Queue)

(21)

In a D-Queue both insertion and deletion operations are performed at either end of the queue. i.e., we performed insertion and deletion from rear as well as front end.



Types of D-Queue \rightarrow

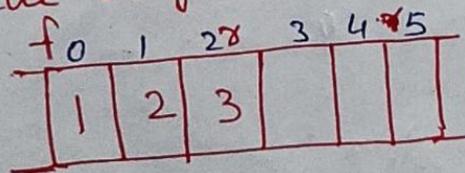
- (i) Input restricted D Queue \rightarrow Element can be added at only one end but we can delete the element from both end.
- (ii) Output restricted D Queue \rightarrow Element can be deleted only from one end but allow insertion at both ends.

Operations on D-Queue \rightarrow (Array implementation)

(1) Insertion at front \rightarrow

Concept \rightarrow

Case 1 \rightarrow Queue is full at front

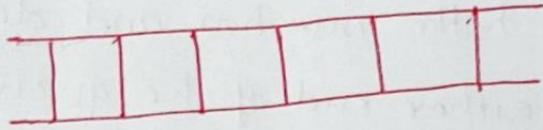


if (front == 0)

printf("Overflow");

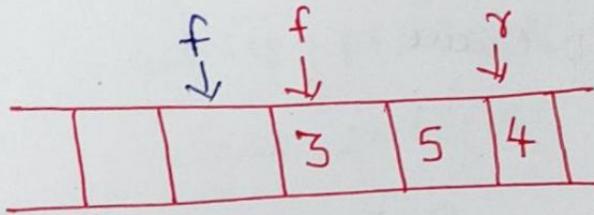
Case 2 → Queue is empty

f = -1
r = -1



if (front == -1 or rear == -1)
front = rear = 0

Case 3:- Queue is neither empty nor full at front



front = front - 1;
DQ[front] = item

Algorithm →

Insert_at_front ()

1. Read item
2. if (front == 0) then
write "Overflow" & exit
3. if (front == -1)
front = rear = 0
else
front = front - 1
4. DQ[front] = item
5. Exit